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Jean-François Raskin\*

**Nicolas Mazzocchi\***

Sriram Sankaranarayanan†

Ashutosh Trivedi†

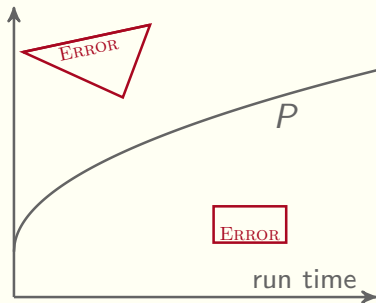
\* Université libre de Bruxelles

† University Colorado Boulder

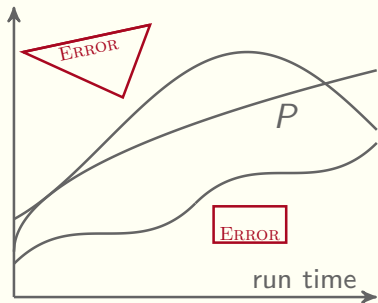
*CONCUR 2020*

# Weighted Transducers for Robustness Verification

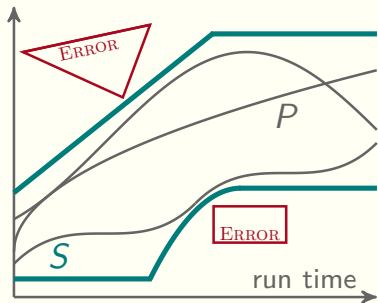
# Model-checking



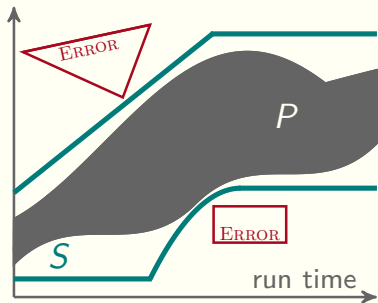
# Model-checking



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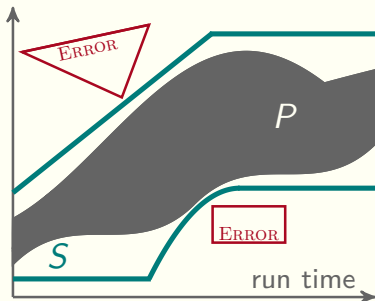


# Model-checking

**P**rogram  
traces

**S**pecified  
traces

$P \subseteq S$

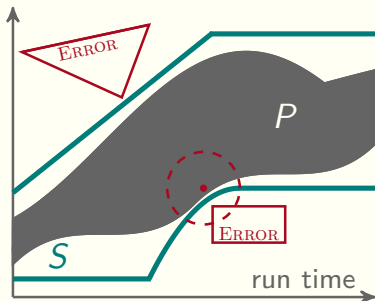


# Model-checking

**P**rogram  
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$$P \subseteq S$$

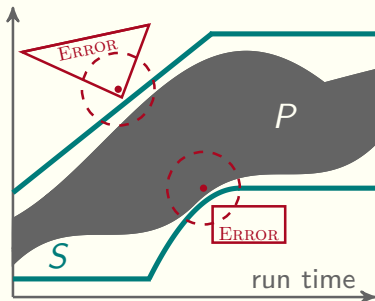


# Model-checking

**P**rogram  
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$P \subseteq S$





# Model-checking

**P**rogram  
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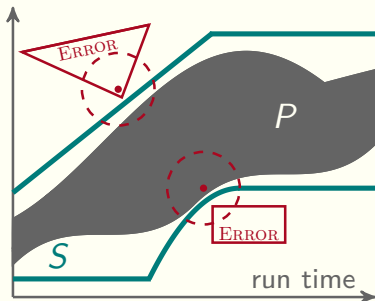
$$P \subseteq S$$

$$P \subseteq_{N,\nu} S$$

**N**oise

threshold  $\nu$

$\nu$ -Neighborhood



# Model-checking

**P**rogram  
traces

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traces

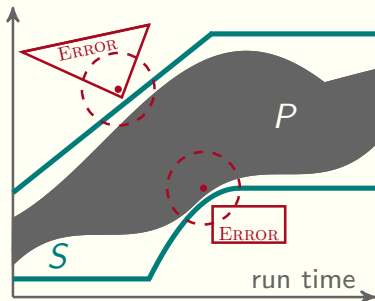
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## Language inclusion problem

- ▶ **Classical**  $\forall t \ t \in P \implies t \in S$

# Model-checking

**P**rogram  
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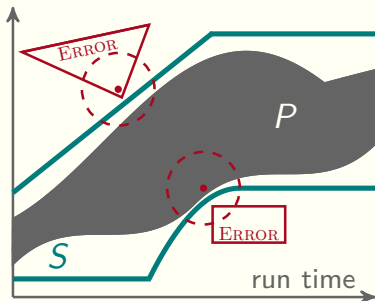
$$P \subseteq S$$

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**N**oise

threshold  $\nu$

$\nu$ -Neighborhood



## Language inclusion problem

- ▶ **Classical**  $\forall t \ t \in P \implies t \in S$
- ▶  **$\nu$ -Robust**  $\forall t_1, t_2 \ (N(t_1, t_2) \leq \nu \wedge t_1 \in P) \implies t_2 \in S$

# Quantitative formalism

Accepting run



# Quantitative formalism

Accepting run



# Quantitative formalism

## Accepting run



## Noise model

$$u_{in} = a_0 \dots a_n \overset{\text{noise}}{\rightsquigarrow} u_{out} = w_0 \dots w_n$$

$$N(u_{in}, u_{out}) = \text{cost}(x_0, \dots, x_n)$$

$$\text{cost}: \mathbb{N}, \dots, \mathbb{N} \rightarrow \mathbb{Q}_{\geq 0}$$

# Quantitative formalism

## Accepting run



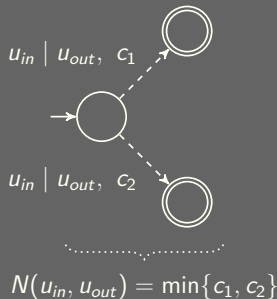
## Noise model

$$u_{in} = a_0 \dots a_n \overset{\text{noise}}{\rightsquigarrow} u_{out} = w_0 \dots w_n$$

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## Non-determinism



# Quantitative formalism

## Accepting run



## Noise model

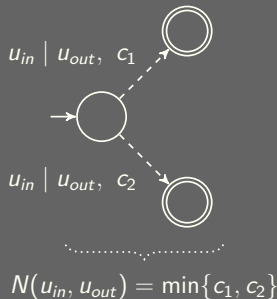
$$u_{in} = a_0 \dots a_n \overset{\text{noise}}{\rightsquigarrow} u_{out} = w_0 \dots w_n$$

$$N(u_{in}, u_{out}) = \text{cost}(x_0, \dots, x_n)$$

$$\text{cost}: \mathbb{N}, \dots, \mathbb{N} \rightarrow \mathbb{Q}_{\geq 0}$$

$$N(u, u) = 0$$

## Non-determinism





# Neighborhood functions

## Edit distance

R E L E   V A N T  
E L E P H A N T

least { inserts  
deletes  
changes

$$\text{Sum} = x_0 + \cdots + x_n$$

Mohri, CIAA'02

# Neighborhood functions

## Edit distance

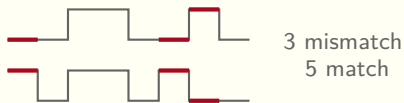
R E L E   V A N T  
E L E P H A N T

least  $\left\{ \begin{array}{l} \text{inserts} \\ \text{deletes} \\ \text{changes} \end{array} \right.$

$$\text{Sum} = x_0 + \dots + x_n$$

Mohri, CIAA'02

## Average mismatch



$$\text{Avg} = \frac{x_0 + \dots + x_n}{n+1}$$

# Neighborhood functions

## Edit distance

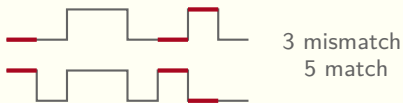
R E L E   V A N T  
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least  $\left\{ \begin{array}{l} \text{inserts} \\ \text{deletes} \\ \text{changes} \end{array} \right.$

$$\text{Sum} = x_0 + \dots + x_n$$

Mohri, CIAA'02

## Average mismatch



$$\text{Avg} = \frac{x_0 + \dots + x_n}{n+1}$$

## Discounted-sum distance

1	0	0	1	0	0	0	$\equiv 72$
0	1	1	1	1	1	0	$\equiv 66$

$$\text{Disc} = \lambda^0 x_0 + \dots + \lambda^n x_n$$

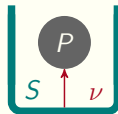
$$\lambda \in \mathbb{Q} \cap (0, 1)$$

# Contributions

## Robust inclusion

**Input:**  $P, S, N, \nu$

**Output:** Decide whether  $P \subseteq_{N, \nu} S$

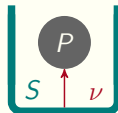


# Contributions

## Robust inclusion

Input:  $P, S, N, \nu$

Output: Decide whether  $P \subseteq_{N, \nu} S$



## Monotonicity

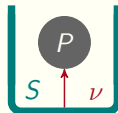
$$\nu_2 > \nu_1 \implies P \subseteq_{N, \nu_2} S \implies P \subseteq_{N, \nu_1} S$$

# Contributions

## Robust inclusion

Input:  $P, S, N, \nu$

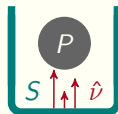
Output: Decide whether  $P \subseteq_{N, \nu} S$



## Threshold synthesis

Input:  $P, S, N$

Output: Compute  $\hat{\nu}$  such that  $P \subseteq_{N, \hat{\nu}} S$



## Monotonicity

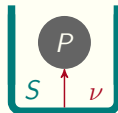
$$\nu_2 > \nu_1 \implies P \subseteq_{N, \nu_2} S \implies P \subseteq_{N, \nu_1} S$$

# Contributions

## Robust inclusion

Input:  $P, S, N, \nu$

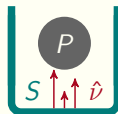
Output: Decide whether  $P \subseteq_{N, \nu} S$



## Threshold synthesis

Input:  $P, S, N$

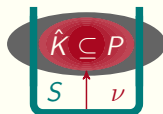
Output: Compute  $\hat{\nu}$  such that  $P \subseteq_{N, \hat{\nu}} S$



## Kernel synthesis

Input:  $P, S, N, \nu$

Output: Compute  $\hat{K} \subseteq P$  such that  $\hat{K} \subseteq_{N, \nu} S$



# Robust inclusion



# Towards an algorithm

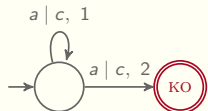
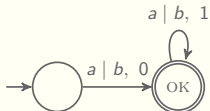
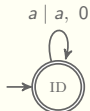
## Example

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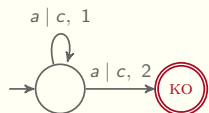
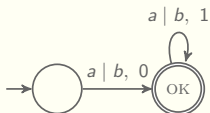
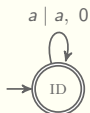
# Towards an algorithm

## Example



# Towards an algorithm

## Example

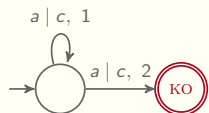
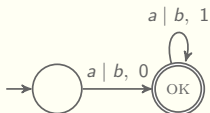
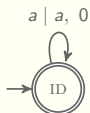


→ 2

	$P \subseteq_{N,\nu} S$
Sum	
Avg	
Disc	

# Towards an algorithm

## Example



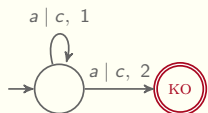
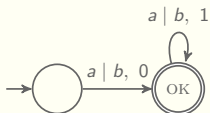
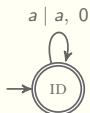
$\rightarrow 2$

$\curvearrowright \rightarrow 3$

	$P \subseteq_{N, \nu} S$
Sum	
Avg	
Disc	

# Towards an algorithm

## Example



$\rightarrow 2$

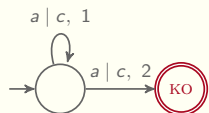
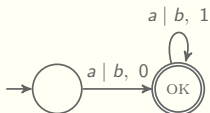
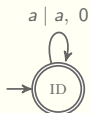
$\hookrightarrow \rightarrow 3$

$\hookrightarrow \dots \hookrightarrow \rightarrow +\infty$

	$P \subseteq_{N, \nu} S$
Sum	
Avg	
Disc	

# Towards an algorithm

## Example



$\rightarrow 2$

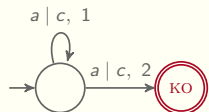
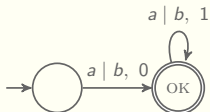
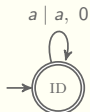
$\hookrightarrow \rightarrow 3$

$\hookrightarrow \dots \hookrightarrow \rightarrow +\infty$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	
Disc	

# Towards an algorithm

## Example

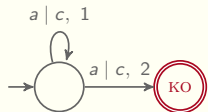
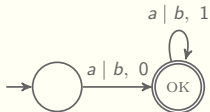
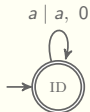


→ 2

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	
Disc	

# Towards an algorithm

## Example



$\rightarrow 2$

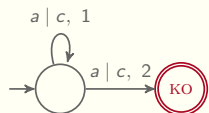
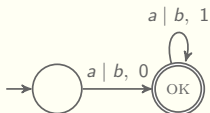
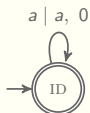
$\curvearrowright \rightarrow 1.5$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	
Disc	



# Towards an algorithm

## Example



$\rightarrow 2$

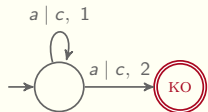
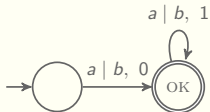
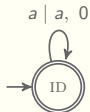
$\hookrightarrow \rightarrow 1.5$

$\hookrightarrow \dots \hookrightarrow \lim_{n \rightarrow +\infty} \frac{n+1}{n}$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	
Disc	

# Towards an algorithm

## Example



→ 2

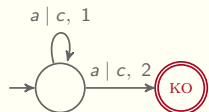
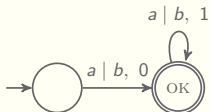
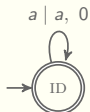
↪ → 1.5

↪...↪ →  $\lim_{n \rightarrow +\infty} \frac{n+1}{n}$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	$\nu \in [0, 1]$
Disc	

# Towards an algorithm

## Example

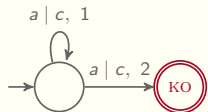
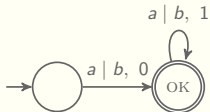
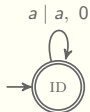


→ 2

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	$\nu \in [0, 1]$
Disc	

# Towards an algorithm

## Example



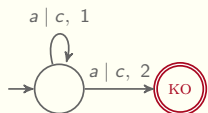
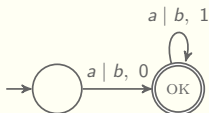
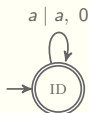
$$\longrightarrow 2$$

$$\text{⤷} \longrightarrow 1 + 2\lambda$$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2)$
Avg	$\nu \in [0, 1]$
Disc	

# Towards an algorithm

## Example



$$\rightarrow 2$$

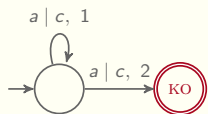
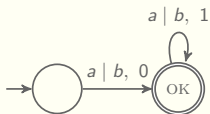
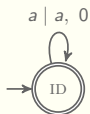
$$\curvearrowright \rightarrow 1 + 2\lambda$$

$$\curvearrowright \dots \curvearrowright \rightarrow \lim_{n \rightarrow +\infty} \frac{1 - \lambda^{n+1}}{1 - \lambda} + 2\lambda^{n+1}$$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2]$
Avg	$\nu \in [0, 1]$
Disc	

# Towards an algorithm

## Example



$$\rightarrow 2$$

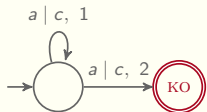
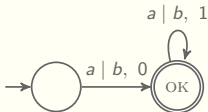
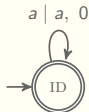
$$\curvearrowright \rightarrow 1 + 2\lambda$$

$$\curvearrowright \dots \curvearrowright \rightarrow \lim_{n \rightarrow +\infty} \frac{1 - \lambda^{n+1}}{1 - \lambda} + 2\lambda^{n+1}$$

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2]$
Avg	$\nu \in [0, 1]$
Disc	$\nu \in [0, \frac{1}{1-\lambda}]$

# Towards an algorithm

## Example

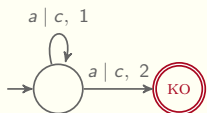
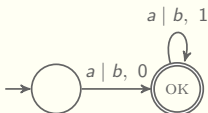
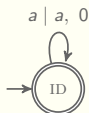


1. identify the **erroneous** outputs automata product
2. compute its smallest value  
Dijkstra, Karp, Alur et al. LATA'13
3. determine its feasibility

	$P \subseteq_{N, \nu} S$
Sum	$\nu \in [0, 2]$
Avg	$\nu \in [0, 1]$
Disc	$\nu \in [0, \frac{1}{1-\lambda}]$

# Towards an algorithm

## Example



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**Robust inclusion:**  $P_{\text{TIME}}$  with  $P, S$  given as NFA, DFA



# Application

## Type-1 Diabetes

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**Hypoglycemia** <70 mg/dl glucose levels for >3h

**Hyperglycemia** >300 mg/dl glucose levels for >3h

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**Noise:** Glucose sensor & Symptoms

# Application

## Type-1 Diabetes

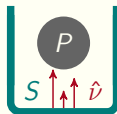
**Hypoglycemia** <70 mg/dl glucose levels for >3h

**Hyperglycemia** >300 mg/dl glucose levels for >3h

**Noise:** Glucose sensor & Symptoms

## Threshold Synthesis

- ▶ **P**atient glucose record
- ▶ **S**ymptoms



# Robust kernel

# Sum operation

## Robust kernel is **regular**

$$\hat{K} = \{u_{in} \in P \mid \forall u_{out} N(u_{in}, u_{out}) \leq \nu \Rightarrow u_{out} \in S\}$$

$$\overline{\hat{K}} = \{u_{in} \in P \mid \exists u_{out} N(u_{in}, u_{out}) \leq \nu \wedge u_{out} \notin S\}$$

$$\overline{\hat{K}} = N_{P, \mathcal{S}}^{\leq \nu}$$

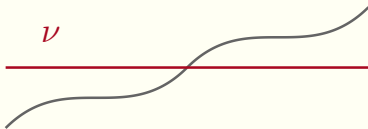
# Sum operation

Robust kernel is **regular**

$$\hat{K} = \{u_{in} \in P \mid \forall u_{out} N(u_{in}, u_{out}) \leq \nu \Rightarrow u_{out} \in S\}$$

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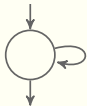
Cost at runtime

Weights in  $\mathbb{N}$   
non-decreasing cost

Kernel emptiness is PSPACE-C

# Avg operation

Robust kernel is **not regular**



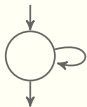
$a \mid a, 0$   
 $b \mid b, 0$   
 $a \mid b, 2$

Let  $\nu = 1$

At most half the input swap  
One  $a$  remain if  $\#(a) > \#(b)$

# Avg operation

Robust kernel is **not regular**



$a \mid a, 0$   
 $b \mid b, 0$   
 $a \mid b, 2$

Let  $\nu = 1$   
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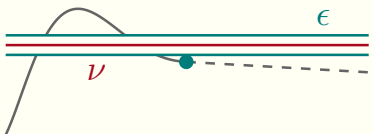
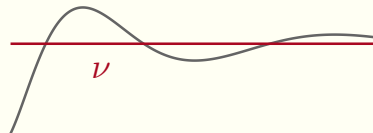
Cost at runtime

Threshold crossed  
arbitrarily many times

Kernel emptiness is undecidable



# Discounted-sum operation



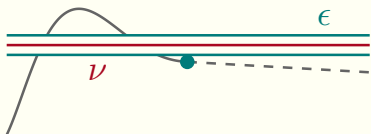
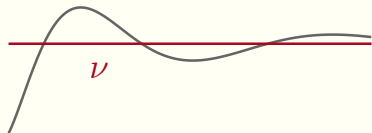
Cost at runtime

## Automata open problem

$$\forall u \quad A_{Disc}(u) \leq \nu$$

Henzinger et al. CSL'08 & LICS'15

# Discounted-sum operation



Cost at runtime

## Automata open problem

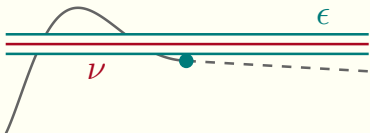
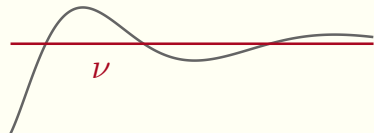
$$\forall u \quad A_{Disc}(u) \leq \nu$$

Henzinger et al. CSL'08 & LICS'15

## $\nu$ -Isolation hypothesis

$$\exists \epsilon \forall u \quad \begin{cases} \nu - \epsilon > A_{Disc}(u) \\ \nu + \epsilon < A_{Disc}(u) \end{cases}$$

# Discounted-sum operation



Cost at runtime

## Automata open problem

$$\forall u \quad A_{Disc}(u) \leq \nu$$

Henzinger et al. CSL'08 & LICS'15

## $\nu$ -Isolation hypothesis

$$\exists \epsilon \forall u \quad \begin{cases} \nu - \epsilon > A_{Disc}(u) \\ \nu + \epsilon < A_{Disc}(u) \end{cases}$$

## Regular $\nu$ -isolated Kernel

$$\overline{\hat{K}} = N_{P, \beta}^{\leq \nu}$$

Kernel emptiness is decidable

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# Weighted Transducers for Robustness Verification