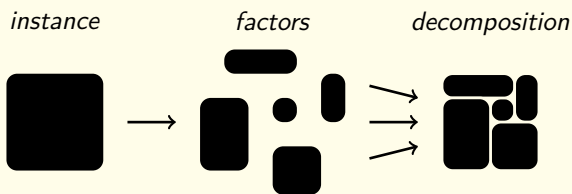


Decomposing Permutation Automata

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Compositionality



Hardware application : Simplify design

Verification application : Rewrite specification

Formalization for DFAs

- The DFA \mathcal{B} is a **factor** of \mathcal{A} if:

$$|\mathcal{B}| < |\mathcal{A}| \quad \wedge \quad L(\mathcal{A}) \subseteq L(\mathcal{B})$$

- \mathcal{A} is a **k -composite** if:

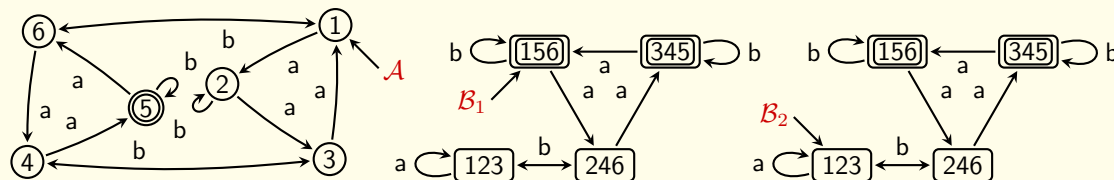
$$L(\mathcal{A}) = \bigcap_{i=1}^k L(\mathcal{B}_i) \text{ where each } \mathcal{B}_i \text{ is a factor}$$

- \mathcal{A} is **composite** if it is k -composite for some k , otherwise it is prime.

- The **orbit** DFA of \mathcal{A} induced by the subset of states S is defined as the DFA obtain by subset construction on \mathcal{A} , where S is the initial state

Key Approach

- A permutation DFA \mathcal{A} is composite iff it can be decomposed into polynomially many orbit DFAs
- If \mathcal{A} is commutative, orbit DFAs defined with a state-space that partitions the state-space of \mathcal{A} suffice



Permutation DFA decomposable into the orbit DFAs induced by $\{1, 5, 6\}$ and $\{1, 2, 3\}$: $L(\mathcal{A}) = \bigcap_{i=1}^2 L(\mathcal{B}_i)$

Summary

Automata class	Composite?	k -Composite?
DFAs	EXPSpace [1]	PSPACE
Permut. DFAs	NP/FPT*	PSPACE
Commut. permut. DFAs	NLOGSPACE	NP-COMPLETE
Unary DFAs	LOGSPACE [2]	LOGSPACE

* Fixed Parameter Tractable in the number of rejecting states

Large decomposition

For infinitely many $n, m \in \mathbb{N}$, there exist a commutative permutation DFA with n states and m letters, requiring $(\sqrt[m]{n} - 1)^{m-1}$ factors to be decomposed.



O. Kupferman and J. Mosheiff
Prime languages
 In J. of Inf. and Comput. 2015



I. Jecker, O. Kupferman and N. Mazzocchi
Unary Prime languages
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